

Stochastic Resource Allocation

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Stochastic Resource Allocation

The Problem: Fixed Resource Bundles

- Resources in the cloud are underutilized
- ► The main cause of resource underutilization is fixed performance bundles



- Clients rent the resources to sustain their highest workload
 - But they do not use the resources all the time
- The provider guarantees with good probability that the clients will be able to use their rented resources at any given time
- It must reserve these resources
 - It cannot resell them or use them to other purposes
- Incentivizing clients to reduce their fixed reserved resource requirements might solve the problem by allowing more clients per physical machine



The existing solution is not perfect



- Still reserve the resource
- Instead of an upper limit...
- Limits the client to a certain average resource consumption
- Adopted by many major cloud providers



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Disadvantages:

- ► Hidden information regarding resource availability
- **Coupling** of reserved resources and average usage

Stochastic Allocation (SA)



- Under the SA mechanism, the provider offers clients a combination:
 - an amount of reserved resources
 - with a choice of a stochastic allocation class
- The provider posts fixed unit-prices for both goods
- And periodically publishes statistics on resource availability for each SA class
- Each client may choose to rent reserved and/or stochastic resources

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- CFS does not support a key feature of SA: defining a different consumption share for the leftover CPUs



- We adapted CFS to support asymmetric reserved resources and share allocations
- We duplicated the CFS logic, to have a second, alternative, CFS

Evaluating Our Solution

How can we compare our solution to **burstable performance** and **fixed performance**?



Will it improve the **utilization**?



Will it be more **popular among clients**?



Will it be more profitable to the providers?

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Will it be more profitable to the providers?

- ► We developed a framework to evaluate new resource allocation schemes
- It simulates a realistic data center with realistic servers and clients

Azure Public Dataset

► We used data from the Azure public dataset



- Includes data for over 2 million clients
 - Purchased bundle (fixed-performance)
 - CPU usage every 5 minutes (min, max and average)
 - and more...

Available from: https://github.com/Azure/AzurePublicDataset



We simulated scenarios where clients share CPU at fine granularity



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▶ We simulated scenarios where clients share CPU at fine granularity



Clients' Load

- We simulated scenarios where clients share CPU at fine granularity
 - We generated 25 samples from each 5 minute sample such that their minimum, maximum and average match the sample
 - We used beta distribution, which can be defined by its average and bounds





What is the performance the clients gain from the CPU?

- To allow the clients to make an informed decision when selecting a bundle, we generated a required performance distribution function
 - Cumulative distribution function
 - It is inspired by performance functions for real applications













- Each simulated client selected the most profitable bundle for its load and resource requirements
- It used its own load statistics to make a decision





- ► To allocate clients to 64-core servers, we randomly shuffled them
- Then, one at a time, each client was assigned to the first server that could accommodate its bundle





- Each client's load for the current day (iteration) was selected cyclically from its data over multiple days
- ► The provider collected statistics on the resource utilization in each server





The cloud provider supplies statistical information regarding the maximal resource amount that a client might obtain over a short period with the commensurate number of shares







- A number of clients were allowed to switch their bundle in each iteration
- They used their own load statistics and the provider's statistical description of the resources that every bundle yields







Evaluating Our Framework

- We simulated a fixed-performance allocation scheme
- Our results were similar to known cloud data before burstable performance was introduced
- ▶ 15%-20% CPU utilization
- Bundle distribution (right)
 - The selected number of virtual cores in our simulation and in the Azure dataset



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Our framework is validated and consistent with real data

Evaluating our Solution







- Fixed Performance (FP)
 - ▶ FP always offered to the clients as an alternative
 - A CPU unit costs \$1
- Burstable Performance (BP) share costs \$2 to \$4
 - The client can rent bundles in which the number of reserved resources equals the number of shares
 - The shares can be utilized without limitations
 - The client's average consumption is limited
- ► Stochastic Allocation (SA) share costs \$0.15 to \$0.9
 - ► The client can rent shares alongside reserved resources
 - A share can be utilized only up to its absolute value

Clients per Server (CPS)



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70% more clients per server compared to BP

Clients per Server (CPS)



- 70% more clients per server compared to BP
- 92%-99% of the clients preferred SA over FP
- 56% of the clients preferred SA over BP and FP

Utilization



Utilization



 SA mean total utilization (73%) is higher than for BP (44%)

Utilization



- SA mean total utilization (73%) is higher than for BP (44%)
- BP reserved utilization is similar to its total utilization

Provider Goals



Public cloud providers:



Is it possible that Amazon has lost money by introducing burstable performance?

Introducing Amazon EC2 T3 Instances

Posted On: Aug 21, 2018

Amazon Web Services (AWS) is introducing the next generation Amazon Elastic Compute Cloud (EC2) burstable general-purpose instances, T3. T3 instances offer a balance of compute, memory, and network resources and are designed to provide a baseline level of CPU performance with the ability to burst above the baseline when needed. T3 instances are powered by the AWS Nitro System which includes a lightweight

Provider Goals



Public cloud providers:

- Maximize their profit from renting their machines
- Take servers' operational costs into account



Is it possible that Amazon has lost money by introducing burstable performance?

Probably not...

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 SA can increase the profits of the public cloud provider by over 28% compared to BP





Public cloud providers:

- Maximize their profit from renting their machines
- Take servers' operational costs into account



Private cloud providers:

Maximize the aggregated benefit all their clients draw from a single server

Provider's Aggregated Benefit per Server



- SA increases the value each server generates for the corporation by over 55% compared to BP
- SA achieved over 98% of the optimal social welfare

\bigcirc Conclusions



- Stochastic CPU allocation via shares allows clients to reduce their reserved resource requirements
- ► SA increases the number of clients per server by more than 70% compared to BP
- ► SA increases the profits of the public cloud provider by over 28% compared to BP
- ▶ SA increases the value each server generates for the corporation by over 55%



- Our evaluation framework is validated against real cloud data
- It is available as an open source: https://bitbucket.org/funaro/stochastic-allocation

Questions?

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imulating Clients



- Real (human) clients may choose an offering in any way they like
- They may choose randomly, take some time to make a decision, or go through a long iterative process of selection and improvement



In the simulation we needed to create realistic artificial intelligence agents which mimic the behavior of real clients

Elients' Requirements

Long-term Requirements

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- ► Have **non-interactive** workloads
- Value finishing the workload by or before a deadline
- Not value getting partial results ahead of time

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Immediate Requirements

- Runs brief independent workloads or an interactive workload, and sleeps the rest of the time
- The failure or fulfillment of one workload does not affect the client's future requirements

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Example: Web Application

- Partition their budget proportionately to the gain from satisfying these dual requirements
- Would not like to miss an opportunity to show an advertisement to their visitors
- Preserve their customers' visit rate



We used the client's purchased fixed-performance bundle as an lower bound on their budget and draw a budget out of a Pareto distribution that is higher than this bound



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- We used the client's purchased fixed-performance bundle as an lower bound on their budget and draw a budget out of a Pareto distribution that is higher than this bound
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- A probability density function (PDF) of the portion between the two valuation types
- The immediate requirements portion completes it to 1



Valuation Functions

- For each type of requirement, we created a valuation function
- V_{imm} and V_{lt} yields a monetary value for an **immediate** and **long-term** performance
- ▶ The client's profit: $E(V_{imm}(P_{r,s})) + V_{lt}(E(P_{r,s})) Cost_{r,s}$
 - s and r denotes the number of shares and reserved resources
 - \blacktriangleright $P_{r,s}$ denotes a random variable which describes the clients' performance with r and s
- We generated for each client V_{imm} and V_{lt} such that its profit will be higher for their bundle of reserved resources than any other bundle of reserved resources (s ≡ 0)
- Solved implicitly



Comparison of the Social Welfare



Distribution of burstable performance bundle



Distribution of unlimited shares bundle



Distribution of limited shares bundle



Effective vs. expected Clients' Revenue



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- Over the course of the last 60 iterations, up to 12% of the clients changed their selected bundle from the first iteration to the last, in all the tested cases.
- The standard deviation of the selected bundles' distribution over these iterations was under 0.6% and the standard deviation of the shares CDF was under 0.01%, in all the tested cases.



- When we modified the **beta density** to be 0.5, 10 or 50, CPS was increased by up to 6% for SA, on the one hand, and reduced by up to 5% for BP, on the other, compared with the main value (1).
- When we avoided the over-the-top extrapolation of the generated load values, CPS was reduced by up to 7%.
- When we modified the performance functions so they were linear and concave, CPS was reduced by up to 3% compared with monotonically increasing ones.
- ▶ When we modified the **Pareto index**, CPS was reduced by up to 6% for a Pareto index of 0.8 and increased by up to 1% for an index of 1.3, compared with the main index (1.1).



- We also modified the number of clients that can **change their bundle**.
- The average CPS was not affected when 384 (or less) clients changed their bundle at once.
- ▶ When more than 256 clients changed their bundle, however, the results fluctuated.
- ▶ When more than 384 clients changed their bundle, the results failed to converge.
- Throughout the above-mentioned modifications, this ratio turned out higher than in the main results presented earlier.